On the Computational Strength of Calculi for (Bio)Chemistry

Based on joint works with:

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done within the UniBo Strategic Project CompReNDe

(Compositional and executable Representations of Nano Devices)

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A
$$\rightarrow^{r} C_{1}+...+C_{n}$$
 A ::= $\tau@r;C_{1}|...|C_{n}+b@s;\mathbf{0}$
A+B $\rightarrow^{s} D_{1}+...+D_{m}$ B ::= $\overline{b}@s;D_{1}|...|D_{m}$

What is the computational power of (Bio)Chemistry?

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Plan of the talk

- Chemistry as a Computational Model
 - ... not a new issue
- Chemical Ground Form (CGF) [Cardelli08]
 - ... a new way to analyze chemistry
- Computational Strength of Chemistry [Concur08]
 - Non Turing universal but...
 - Turing complete formalisms can be approximated with any degree of precision
- From Chemistry to...
 - Biochemistry: BGF [AB08]
 - Nano devices: nano-K [CMSB07,TCS08]
- Computational strength of BGF and K-calculi [CMSB09]
- Concluding remarks

Is Chemistry Turing powerful?

- Magnasco. Chemical Kinetics is Turing Universal. Phys Rev Lett. 1997
 - Answer: YES... but justification not convincing (only Digital Computers with bounded memory are considered)
- Liekens and Fernando. Turing Complete Catalytic
 Particle Computers. In Proc. ECAL'07. 2007
 - Answer: YES... but justification not convincing (only Minsky Machines with bounded computation are considered)
- Soloveichik et al. Computation with Finite Stochastic Chemical Reaction Networks. Nat. Computing. 2008
 - Answer: NO... but all Minsky and Turing Machines can be at least approximated with any given degree of precision

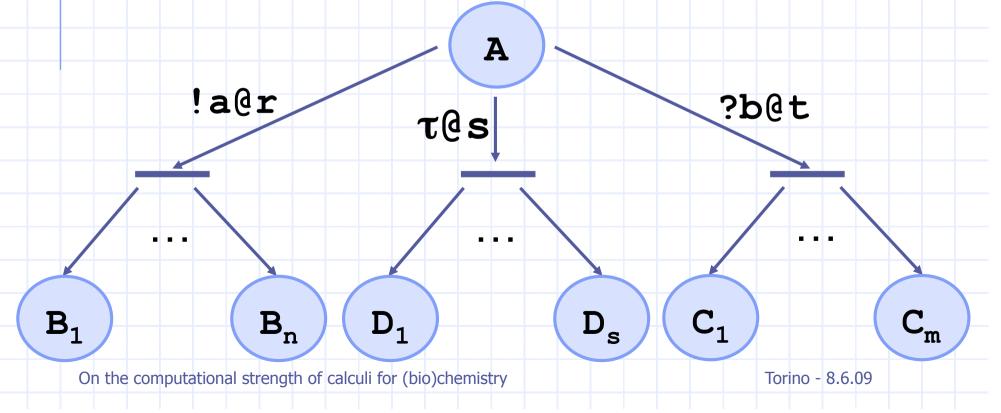


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Chemical Ground Forms

 Stochastic variant of Milner's CCS, with an equivalent graphical notation (Stochastic Collective Automata)



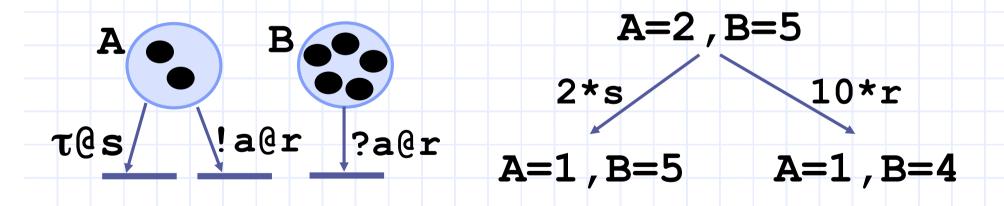
Stochastic semantics

- Actions take (an exponentially distributed amount of) time
 - Internal delay: τ@r
 - Pr(internal delay < t) = 1-e^{-rt}
 - Synchronization between complementary actions: ?a@r, !a@r
 - ◆ Pr(synchronization time < t) = 1-e^{-rt}

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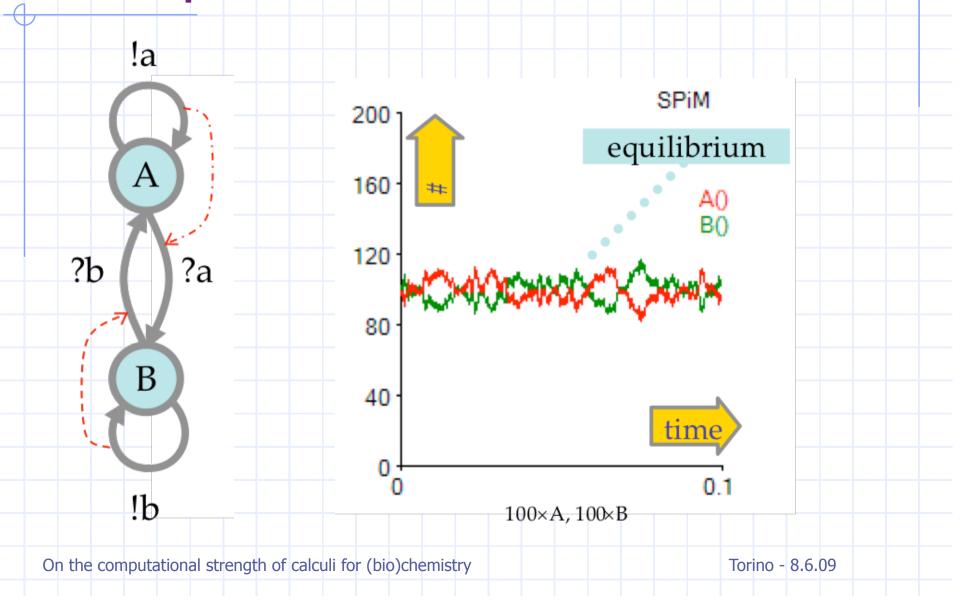
Collective stochastic semantics

 To faithfully model chemistry, the rates of distinct transitions with the same effect must be added

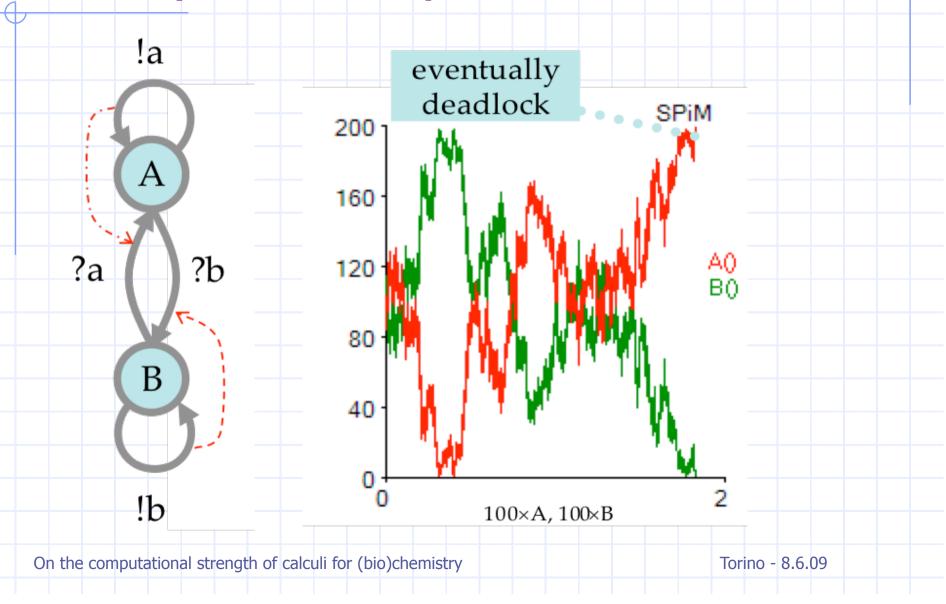


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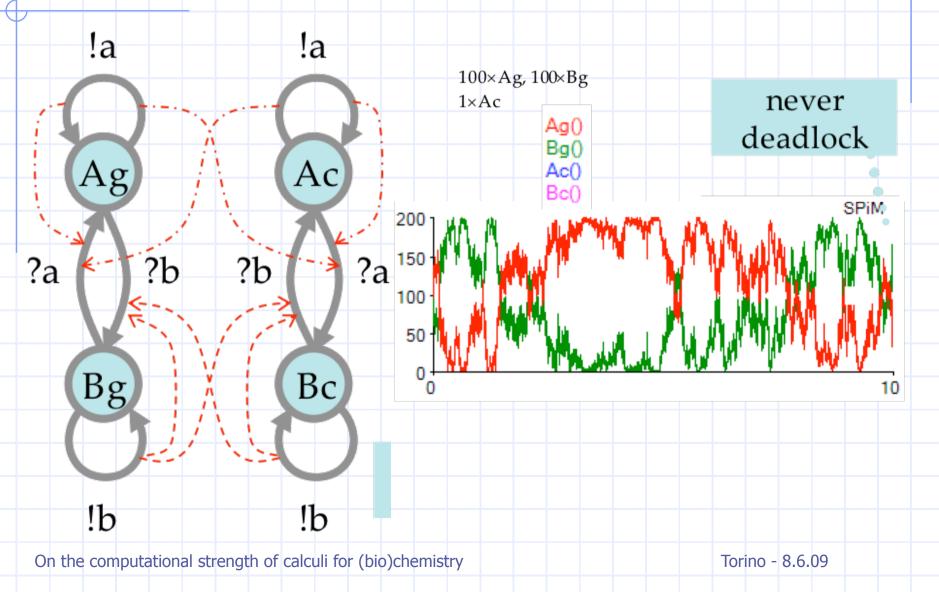
Example: Celebrities...



Example: Groupies...



Example: Groupies with one Celebrity...



CGF = FSCRN

[Cardelli08]

Continuous-State
Semantics

Continuous Chemistry



CGF FSCRN



1

Discrete-State
Semantics

Discrete Chemistry

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Non Turing universality

- The (existential version of) the Halting problem is decidable in CGF
- Proof:
 - We first define an encoding of CGF into (nondeterministic) Petri nets
 - Theorem:
 Given a CGF process P, it can reach a halted configuration (with prob. > 0) iff the net that encodes P has a deadlock
 - Deadlock is decidable in Petri nets

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Approximation of Turing complete formalisms in CGF

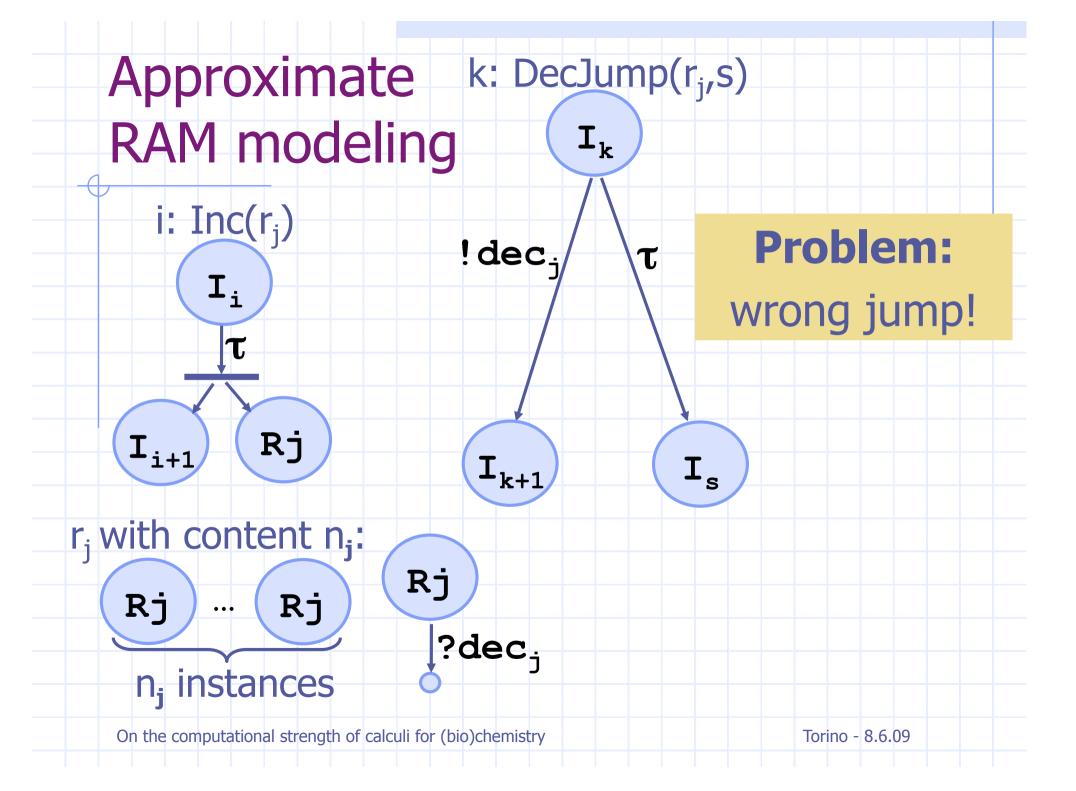
- We approximate (for any degree of precision)
 Random Access Machines (RAMs)
- ◆ RAMs [Min67]:
 - **Registers**: r₁ ... r_n hold natural numbers
 - Program: sequence of indexed instructions
 - i: Inc(r_j): add 1 to the content of r_j and go to the next instruction
 - i: DecJump(r_j,s): if the content of r_j is not 0 then decrease by 1 and go to the next instruction; otherwise jump to instruction s

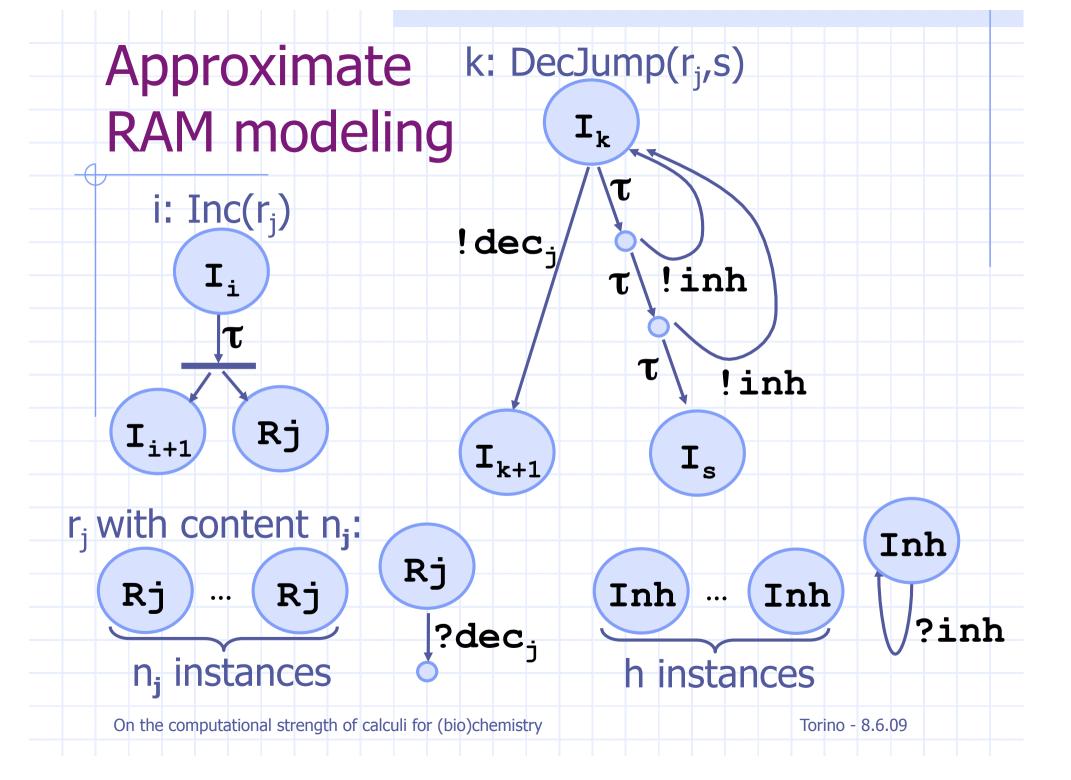
RAM encoding

- RAMs cannot be faithfully modeled by a CGF process
 - otherwise (by decidability of existential term. in CGF)
 RAM termination is decidable

Main theorem:

Given $\varepsilon > 0$ and a RAM, it **can** be modeled by a CGF process that includes also wrong computations, but the probability a wrong computation is scheduled is smaller than ε





Approximate RAM modeling

k: DecJump(r_j,s)

i: $Inc(r_j)$ T_i T_{i+1} R_j

!dec; tinh

 I_{k+1}

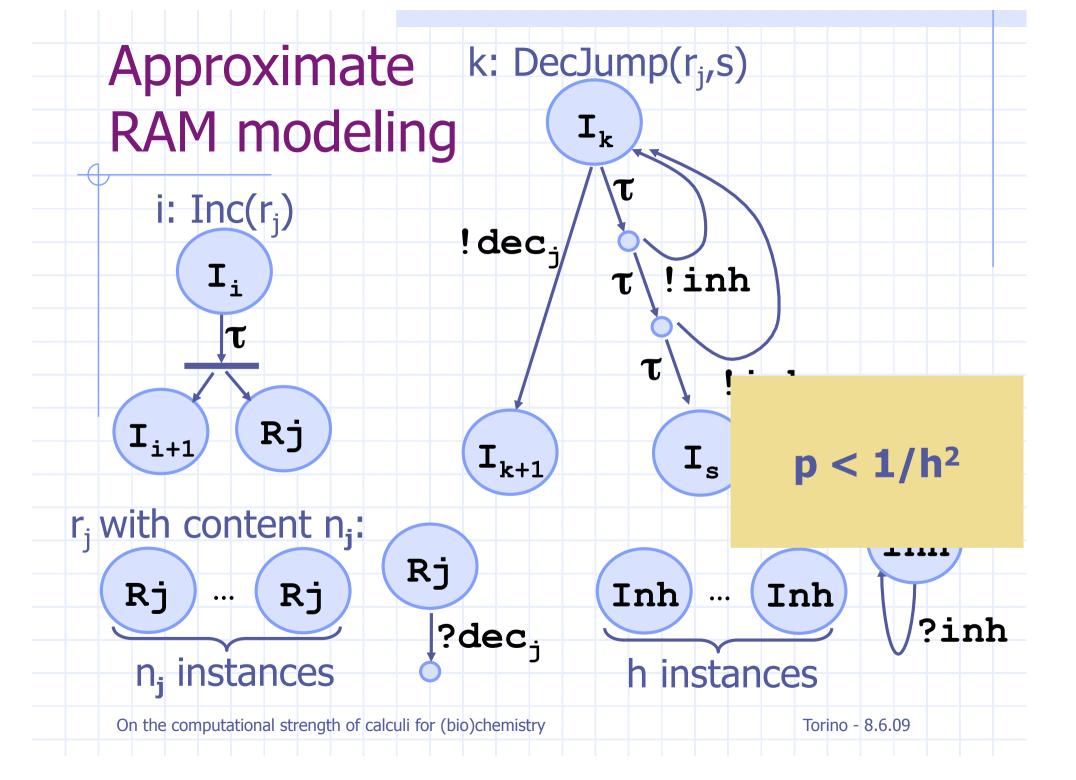
 $1/(n_j+1)*1/h^2 <$ **p < 1/h**²

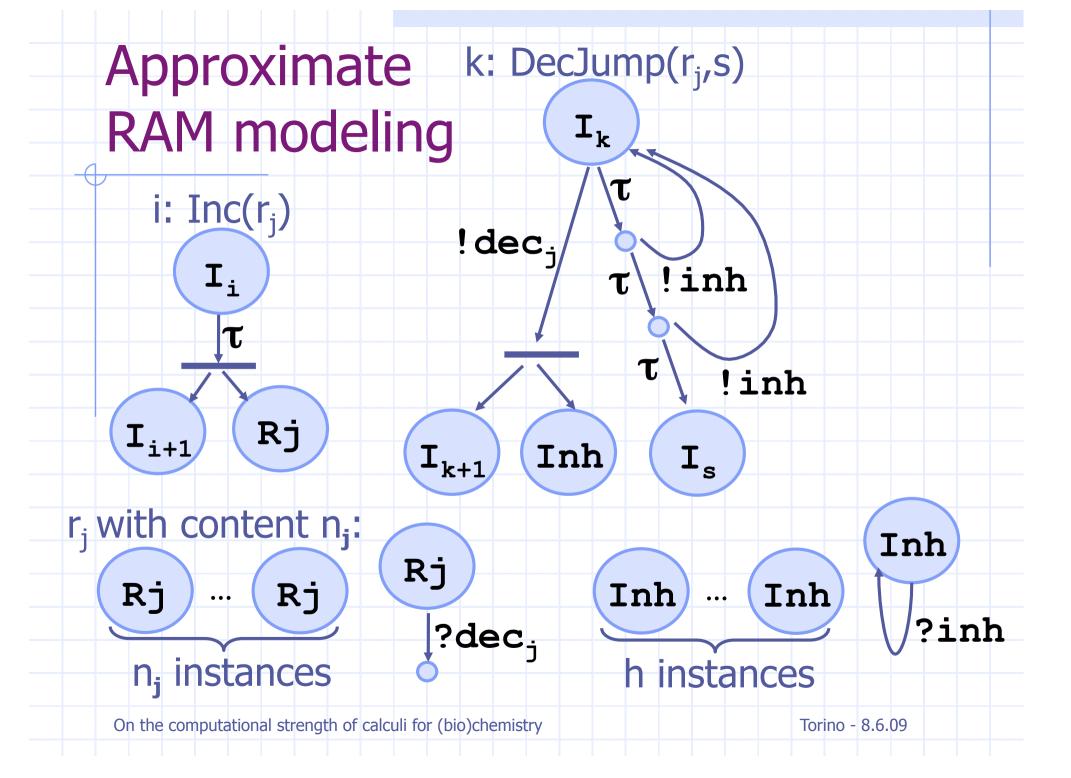
r, with content n.

R

But for unbounded computations, with infinitely many decrements, the prob. of a wrong jump is 1

?inh





k: DecJump(r_i,s) **Approximate** RAM modeling I_k i: Inc(r_i) !dec; !inh τ !inh Rj I_{i+1} I_{k+1} Inh

r_j wi

Incrementing the occurrences of Inh the prob. a wrong jump is taken is

 $p < \sum_{i=inith...\infty} 1/i^2$

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Torino - 8.6.09

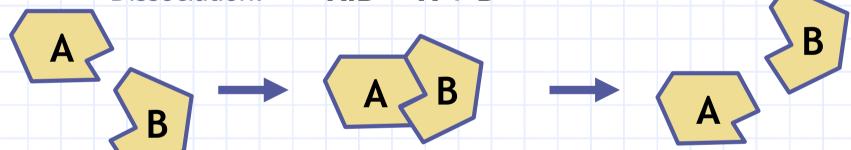
nh

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Basic Biochemistry

- Molecules may also form reversible complexes
 - Association: A + B → A:B
 - Dissociation: A:B → A + B

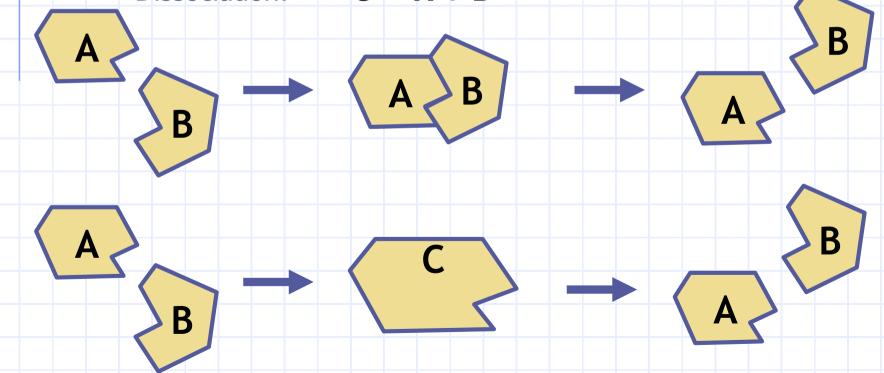


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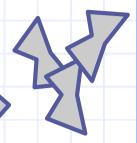
Basic Biochemistry

- Molecules may also form reversible complexes
 - Association: A + B → C
 - Dissociation: $C \rightarrow A + B$

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What's the Difference?



Consider linear polymerization:

The "chemical program" for polymerization:

$$P_0 + M \rightarrow P_1$$

$$P_1 + M \rightarrow P_2$$

$$P_2 + M \rightarrow P_3$$

•••

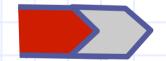
an infinite (non-)program an infinite set of species

$$P_{10757} + M \rightarrow P_{10758}$$

Such specificity is unreal

But "nature's program" for polymerization has to fit e.g. in the genome, so it cannot be infinite! Clearly, nature must be using a different "language" than basic chemistry:





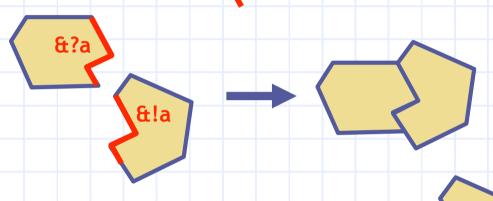
molecule with convex patch +
molecule with concave patch →
molecule with convex patch

Finite program composed of local rules

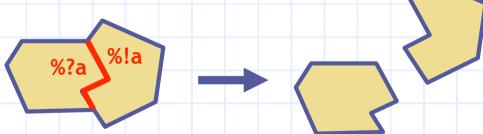
Biochemical Ground Form= CGF + Association and Dissociation

- Association patches are named
- the a shape

- ◆ & association
 - &?a associate
 - &!a co-associate



- % dissociation
 - %?a dissociate
 - %!a co-dissociate



- A given patch can hold only one association at a time
- Two molecules can dissociate only if they are associated

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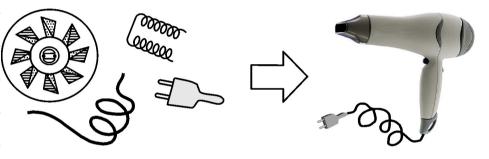
Nano Devices

Supramolecular system

- Built bottom-up: assembling molecular components
- Used to "implement"
 basic functionalities:
 binary memories, basic
 logic gates, mechanical
 movements, ...

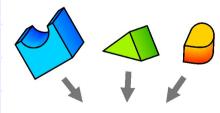
macroscopic components

macroscopic device

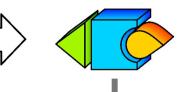


molecular components

molecular device (supramolecular structure)



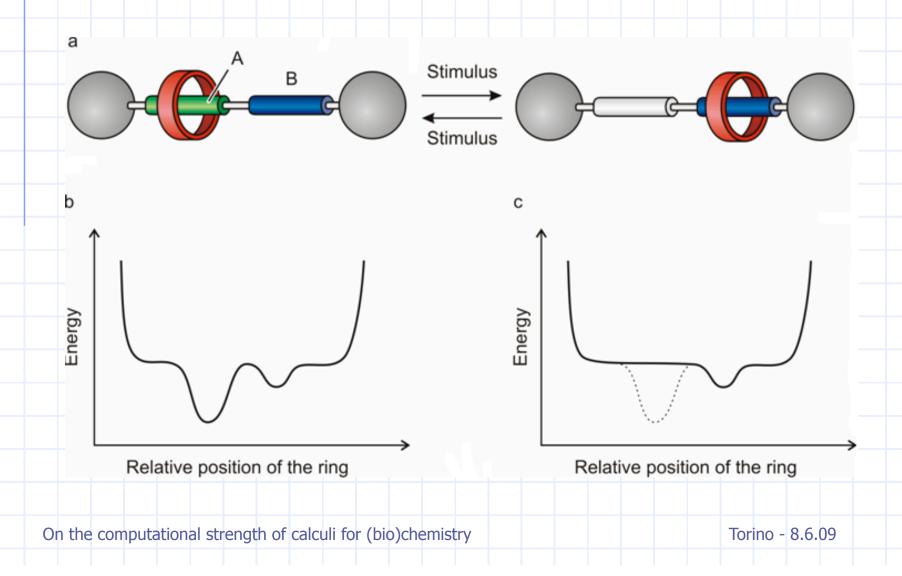




complex function

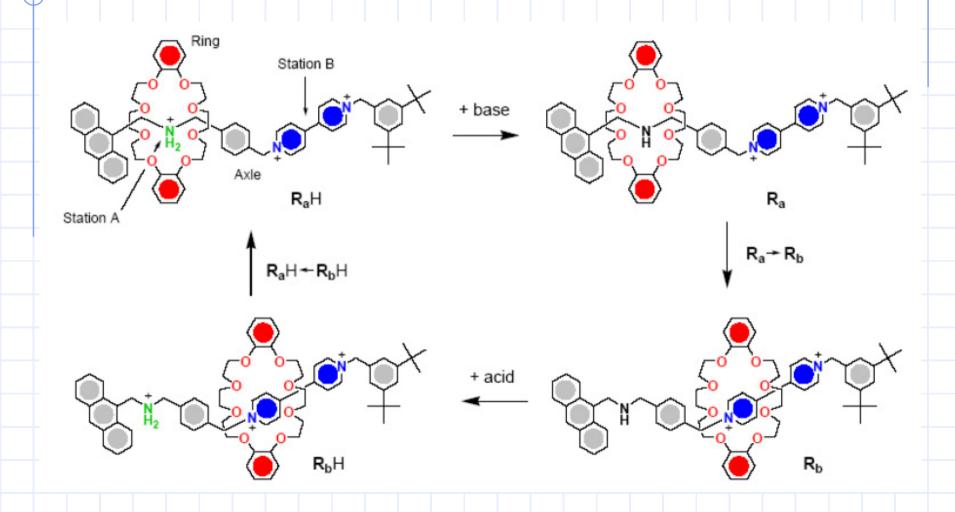
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Example: [2]rotaxane



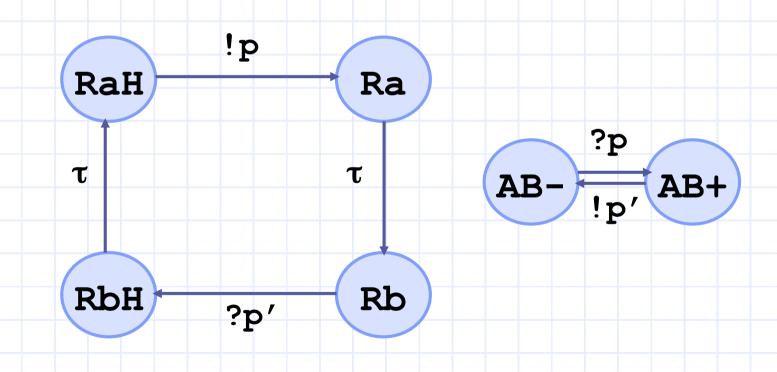
Rotaxane RaH

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Modeling RaH in CGF

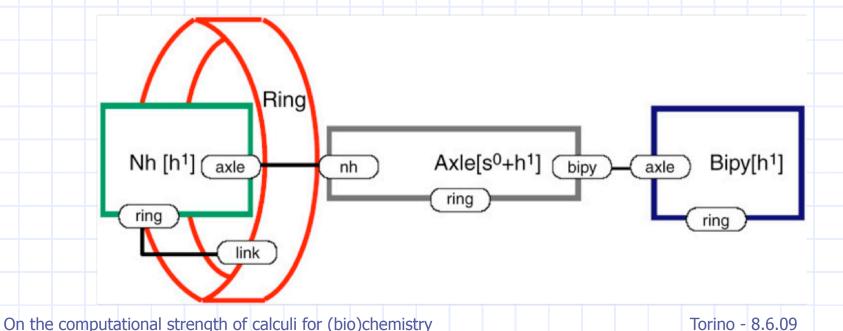
The RaH process calculus definition



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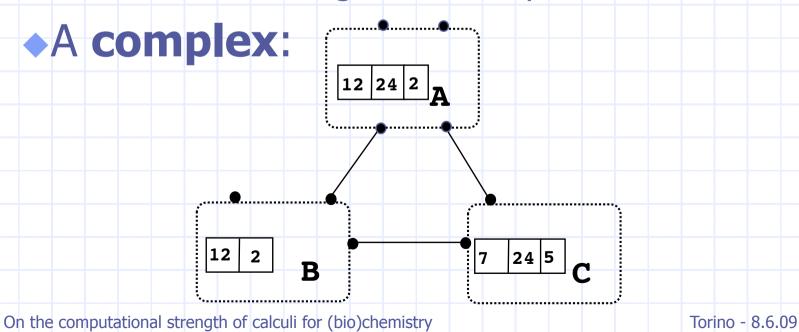
Towards a compositional representation

- In this representation there is no indication of the exploited components!
- ◆ Solution: follow the approach of K-calculi



K-calculi: the nodes

- K is a graph-rewriting system where nodes
 - carry state information (about phosphorilation, shape,...)
 - have sites through which they bind other nodes



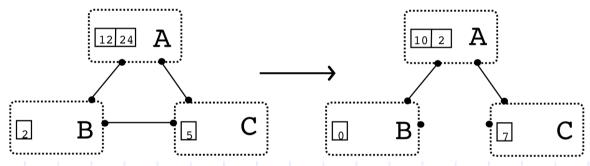
and its syntactic description:

K-calculi: the reactions

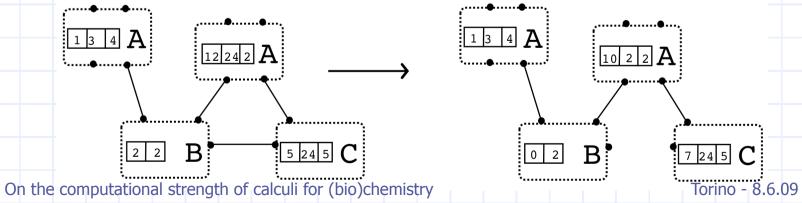
A reaction rule:

 $A[1^{12}+2^{24}](1^{x}+2^{y}), B[1^{2}](2^{x}+3^{z}), C[1^{5}](1^{z}+2^{y})$

$$\rightarrow$$
 A[1¹⁰+2²](1^x+2^y), B[1⁰](2^x+3), C[1⁷](1+2^y)

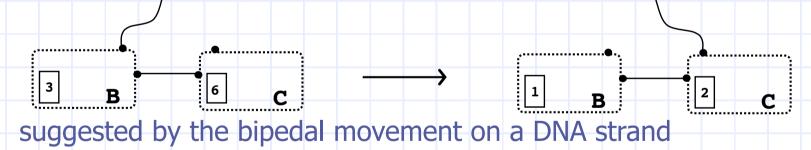


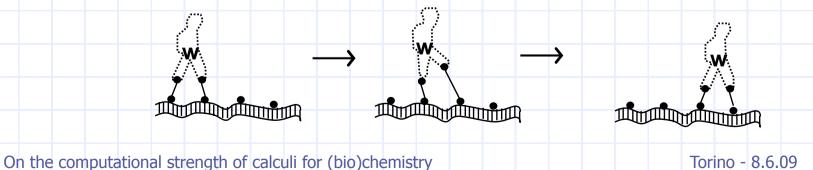
and a possible instantiation:



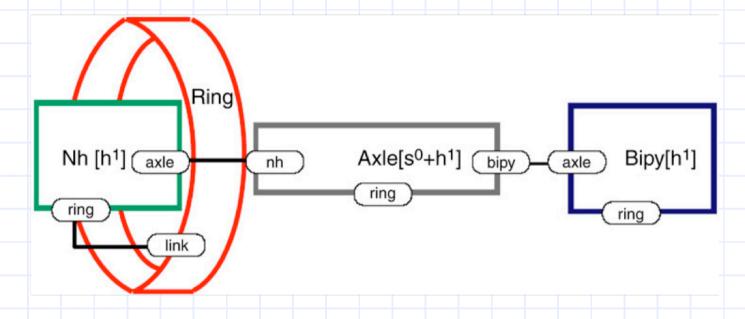
Three types of reactions: creations, destructions and bond flipping

- Reactions are monotone:
 - either creations (bonds and/or proteins are created)
 - or destructions (bonds and/or proteins are removed)
 - or (extension of k) bond flippings (an edge of a bond may be moved from one molecule to another)





RaH modeled in K: the initial state



 $Nh[h^{1}](axle^{s} + ring^{x}),$ $Bipy[h^{1}](axle^{r} + ring),$ $Axle[s^0 + h^1](nh^s + bipy^r + ring),$ $Ring(link^x)$

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RaH modeled in K: some reaction rules

Deprotonation:

$$Nh[h^{1}], Base[h^{0}] \stackrel{nh_base}{\underset{d}{\longleftarrow}} Nh[h^{0}], Base[h^{1}]$$

Disconnection:

 $Nh[h^0](axle^s + ring^x)$, $Axle[s^0](nh^s + ring) \stackrel{\infty}{\rightarrow} Nh[h^0](axle^s + ring)$, $Axle[s^0](nh^s + ring^x)$

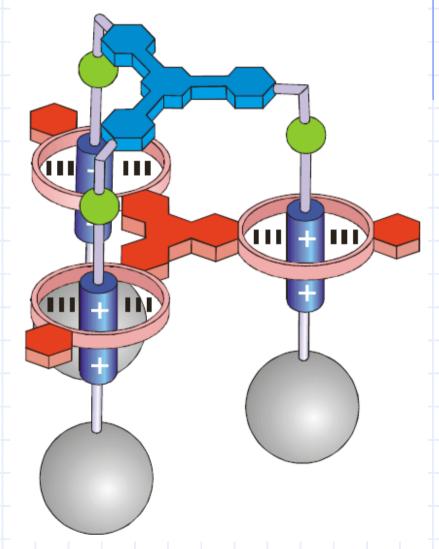
Connection:

$$Axle[s^0](bipy^r + ring^x)$$
, $Bipy[h^0](axle^r + ring)$ $\xrightarrow[unlink_bipy]{link_bipy}$ $Axle[s^1](bipy^r + ring)$, $Bipy[h^0](axle^r + ring^x)$

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From the rotaxane to the nano-elevator

- Nano-elevator:
 three joined rotaxanes
 with a platform connected
 to the rings
- We plan to re-use the modeling of the rotaxane also for the nano-elevator (adding two new species: top & platform)
- Monolithic approach:
 exponential explosion of the species: 4³ = 64



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Turing completeness of BGF

- BGF=CGF+association+dissociation
 - This calculus is Turing complete
- Proof: RAMs can be faithfully modeled representing registers as polymers

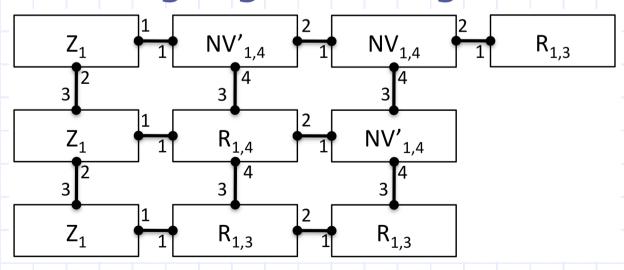
Z_j



 R_{j}

Turing completeness of K-calculi

- K without destructions and internal state is already Turing complete
- Proof: RAMs can be faithfully modeled representing registers as grids



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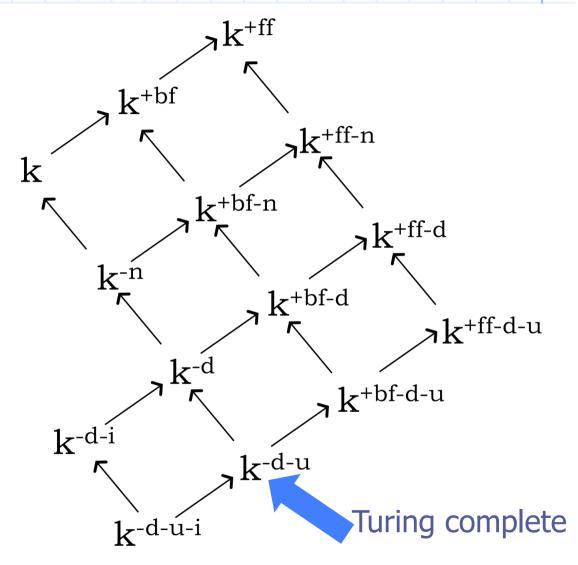
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Conclusion

- CGF, the calculus for chemistry, is not Turing powerful...
 - ...but Turing complete formalisms can be
 approximated with any given degree of precision
- Molecule complexation "Turifies" chemistry
 - BGF includes a basic mechanism for association/ dissociation, which is also biologically realistic, that makes CGF Turing complete

K-calculi: the lattice

- The ordering is the subcalculus relation
- Dialects are inspired by biological phenomena



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K-calculi: the lattice

Reachability

decidable

Reachability/Coverability

undecidable

Coverability (limiting the size of complexes) **decidable**

k

Coverability decidable $k^{-d-\acute{u}-i}$

> k^{+bf-d-u}

 k^{+bf-n}

Turing complete

Coverability undecidable

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Future work

- Understand better the actual importance of molecule dissociation/bond destruction
 - K without bond destruction is Turing complete while BGF without dissocation is not!
 - Possible motivations: multimolecular reactions of K, test of the entire interface in K, ...
- Quantitative analysis
 - The approximated Turing completeness of CGF suggests that quantitative analysis is undecidable (is a molecule/complex reachable with prob. greater than P?)
 - As done for qualitative analysis, we plan to capture fragments
 of these calculi in which also quantitive analysis is possible

References

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